

تمت مشاركة هذه المعلومة بإشارة مشاركة ***أبيض*** حيث يسمح بتبادلها Please note that this notification/advisory has been tagged as TLP ***WHITE*** where information can be shared or published on any public forums.

أو نشرها من خلال القنوات العامة.

في ضوء دور الهيئة الوطنية للأمن السيبراني للمساعدة في حماية الفضاء As part of NCA duties to help securing the cyberspace and protecting national interests, NCA provides the weekly summary of published السيبراني الوطني، تود الهيئة مشاركتكم النشرة الأسبوعية للتُغرات المسجلة vulnerabilities by the National Institute of Standards and Technology the National Institute of Standards and Technology (NIST) من قبل (NIST) National Vulnerability Database (NVD) for the week from 14th ربيو إلى 18 يوليو إلى National Vulnerability Database (NVD) of July to 20th of July. Vulnerabilities are scored using the Common معيار معيار معيار الثغرات باستخدام معيار Vulnerability Scoring System (CVSS) standard as per the following حيث يتم تصنيف الثغرات بناء على Vulnerability Scoring System (CVSS) severity:

Critical: CVSS base score of 9.0-10.0 High: CVSS base score of 7.0-8.9 Medium: CVSS base score 4.0-6.9 Low: CVSS base score 0.0-3.9

التالي:

عالى جدًا: النتيجة الأساسية لـ10.0-CVSS 9.0

عالى: النتيجة الأساسية لـ8.9-7.0 CVSS

متوسط: النتيجة الأساسية لـ6.9-CVSS 4.0

منخفض: النتيجة الأساسية لـ CVSS 0.0-3.9

| CVE ID & Source | Vendor - Product | Description | Publish Date | CVSS Score | Severity |
|-----------------|--|---|--------------|---------------|----------|
| CVE-2024-39736 | ibm - multiple products | IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 is vulnerable to HTTP header injection, caused by improper validation of input by the HOST headers. This could allow an attacker to conduct various attacks against the vulnerable system, including cross-site scripting, cache poisoning or session hijacking. IBM X-Force ID: 296003. | 2024-07-15 | 9.8 | Critical |
| CVE-2024-21181 | oracle - multiple products | Vulnerability in the Oracle WebLogic Server product of Oracle Fusion Middleware (component: Core). Supported versions that are affected are 12.2.1.4.0 and 14.1.1.0.0. Easily exploitable vulnerability allows unauthenticated attacker with network access via T3, IIOP to compromise Oracle WebLogic Server. Successful attacks of this vulnerability can result in takeover of Oracle WebLogic Server. CVSS 3.1 Base Score 9.8 (Confidentiality, Integrity and Availability impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:H/A:H). | 2024-07-16 | 9.8 | Critical |
| CVE-2024-5471 | zohocorp - manageengine_ddi _central | Zohocorp ManageEngine DDI Central versions 4001 and prior were vulnerable to agent takeover vulnerability due to the hard-coded sensitive keys. | 2024-07-17 | 9.8 | Critical |
| CVE-2023-46801 | apache - linkis | In Apache Linkis <= 1.5.0, data source management module, when adding Mysql data source, exists remote code execution vulnerability for java version < 1.8.0_241. The deserialization vulnerability exploited through jrmp can inject malicious files into the server and execute them. This attack requires the attacker to obtain an authorized account from Linkis before it can be carried out. We recommend that users upgrade the java version to >= 1.8.0_241. Or users upgrade Linkis to version 1.6.0. | 2024-07-15 | 8.8 | High |
| CVE-2023-49566 | apache - linkis | In Apache Linkis <=1.5.0, due to the lack of effective filtering of parameters, an attacker configuring malicious db2 parameters in the DataSource Manager Module will result in jndi injection. Therefore, the parameters in the DB2 URL should be blacklisted. This attack requires the attacker to obtain an authorized account from Linkis before it can be carried out. Versions of Apache Linkis <=1.5.0 | 2024-07-15 | 8.8 | High |

| | | ill be offerted | | Г | |
|----------------|---|---|------------|-----|--------|
| | | will be affected. We recommend users upgrade the version of Linkis to version 1.6.0. | | | |
| CVE-2024-3168 | google - chrome | Use after free in DevTools in Google Chrome prior to 122.0.6261.57 allowed a remote attacker to potentially exploit heap corruption via a crafted HTML page. (Chromium security | 2024-07-16 | 8.8 | High |
| CVE-2024-3169 | google - chrome | severity: Medium) Use after free in V8 in Google Chrome prior to 121.0.6167.139 allowed a remote attacker to potentially exploit heap corruption | 2024-07-16 | 8.8 | High |
| CVE-2024-3170 | google - chrome | via a crafted HTML page. (Chromium security severity: High) Use after free in WebRTC in Google Chrome prior to 121.0.6167.85 | 2024-07-16 | 8.8 | High |
| CVL-2024-3170 | google - cilionie | allowed a remote attacker to potentially exploit heap corruption via a crafted HTML page. (Chromium security severity: High) | 2024-07-10 | 0.0 | Tilgii |
| CVE-2024-3171 | google - chrome | Use after free in Accessibility in Google Chrome prior to 122.0.6261.57 allowed a remote attacker who convinced a user to engage in specific UI gestures to potentially exploit heap corruption via specific UI gestures. (Chromium security severity: Medium) | 2024-07-16 | 8.8 | High |
| CVE-2024-3172 | google - chrome | Insufficient data validation in DevTools in Google Chrome prior to 121.0.6167.85 allowed a remote attacker who convinced a user to engage in specific UI gestures to execute arbitrary code via a | 2024-07-16 | 8.8 | High |
| CVE-2024-3173 | google - chrome | crafted HTML page. (Chromium security severity: High) Insufficient data validation in Updater in Google Chrome prior to 120.0.6099.62 allowed a remote attacker to perform OS-level privilege escalation via a malicious file. (Chromium security | 2024-07-16 | 8.8 | High |
| CVE-2024-3174 | google - chrome | severity: High) Inappropriate implementation in V8 in Google Chrome prior to 119.0.6045.105 allowed a remote attacker to potentially exploit object corruption via a crafted HTML page. (Chromium security severity: High) | 2024-07-16 | 8.8 | High |
| CVE-2024-3176 | google - chrome | Out of bounds write in SwiftShader in Google Chrome prior to 117.0.5938.62 allowed a remote attacker to perform an out of bounds memory write via a crafted HTML page. (Chromium | 2024-07-16 | 8.8 | High |
| CVE-2024-39877 | apache - airflow | security severity: High) Apache Airflow 2.4.0, and versions before 2.9.3, has a vulnerability | 2024-07-17 | 8.8 | High |
| | | that allows authenticated DAG authors to craft a doc_md parameter in a way that could execute arbitrary code in the scheduler context, which should be forbidden according to the Airflow Security model. Users should upgrade to version 2.9.3 or later which has removed the vulnerability. | | | |
| CVE-2024-27311 | zohocorp - manageengine_ddi central | Zohocorp ManageEngine DDI Central versions 4001 and prior were vulnerable to directory traversal vulnerability which allows the user to upload new files to the server folder. | 2024-07-17 | 8.8 | High |
| CVE-2022-48834 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: usb: usbtmc: Fix bug in pipe direction for control transfers | 2024-07-16 | 7.8 | High |
| | | The syzbot fuzzer reported a minor bug in the usbtmc driver: | | | |
| | | usb 5-1: BOGUS control dir, pipe 80001e80 doesn't match bRequestType 0 WARNING: CPU: 0 PID: 3813 at drivers/usb/core/urb.c:412 | | | |
| | | usb_submit_urb+0x13a5/0x1970 drivers/usb/core/urb.c:410 Modules linked in: CPU: 0 PID: 3813 Comm: syz-executor122 Not tainted 5.17.0-rc5-syzkaller-00306-g2293be58d6a1 #0 | | | |
| | | Call Trace: <task></task> | | | |
| | | usb_start_wait_urb+0x113/0x530 drivers/usb/core/message.c:58 usb_internal_control_msg drivers/usb/core/message.c:102 [inline] usb_control_msg+0x2a5/0x4b0 drivers/usb/core/message.c:153 | | | |
| | | usbtmc_ioctl_request drivers/usb/class/usbtmc.c:1947 [inline] The problem is that usbtmc_ioctl_request() uses usb_rcvctrlpipe() | | | |
| | | for | | | |
| CVE-2022-48837 | linux - multiple products | all of its transfers, whether they are in or out. It's easy to fix. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.8 | High |
| | | usb: gadget: rndis: prevent integer overflow in rndis_set_response() | | | |
| | | If "BufOffset" is very large the "BufOffset + 8" operation can have an | | | |
| CVE-2022-48847 | linux - multiple | integer overflow. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.8 | High |

| | 1 | watch guous: Fix filter limit check | | | |
|----------------|----------------------|---|------------|-----|------|
| | | watch_queue: Fix filter limit check | | | |
| | | In watch_queue_set_filter(), there are a couple of places where we check | | | |
| | | that the filter type value does not exceed what the type_filter bitmap | | | |
| | | can hold. One place calculates the number of bits by: | | | |
| | | if (tf[i].type >= sizeof(wfilter->type_filter) * 8) | | | |
| | | which is fine, but the second does: | | | |
| | | if (tf[i].type >= sizeof(wfilter->type_filter) * BITS_PER_LONG) | | | |
| | | which is not. This can lead to a couple of out-of-bounds writes due to | | | |
| | | a too-large type: | | | |
| | | (1)set_bit() on wfilter->type_filter (2) Writing more elements in wfilter->filters[] than we allocated. | | | |
| | | Fix this by just using the proper WATCH_TYPENR instead, which is the | | | |
| | | number of types we actually know about. | | | |
| | | The bug may cause an oops looking something like: | | | |
| | | BUG: KASAN: slab-out-of-bounds in watch_queue_set_filter+0x659/0x740 | | | |
| | | Write of size 4 at addr ffff88800d2c66bc by task | | | |
| | | watch_queue_oob/611 | | | |
| | | Call Trace: <task></task> | | | |
| | | dump_stack_lvl+0x45/0x59 print_address_description.constprop.0+0x1f/0x150 | | | |
| | | kasan_report.cold+0x7f/0x11b | | | |
| | | watch_queue_set_filter+0x659/0x740 | | | |
| | | x64_sys_ioctl+0x127/0x190 | | | |
| | | do_syscall_64+0x43/0x90 entry_SYSCALL_64_after_hwframe+0x44/0xae | | | |
| | | Allocated by task 611: | | | |
| | | kasan_save_stack+0x1e/0x40 kasan_kmalloc+0x81/0xa0 | | | |
| | | watch_queue_set_filter+0x23a/0x740 | | | |
| | | x64_sys_ioctl+0x127/0x190 do_syscall_64+0x43/0x90 | | | |
| | | entry_SYSCALL_64_after_hwframe+0x44/0xae | | | |
| | | The buggy address belongs to the object at ffff88800d2c66a0 which belongs to the cache kmalloc-32 of size 32 | | | |
| | | The buggy address is located 28 bytes inside of | | | |
| CVE-2022-48848 | linux - linux_kernel | 32-byte region [ffff88800d2c66a0, ffff88800d2c66c0) In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.8 | High |
| | _ | tracing/osnoise: Do not unregister events twice | | | |
| | | Nicolas reported that using: | | | |
| | | # trace-cmd record -e all -M 10 -p osnoisepoll | | | |
| | | Resulted in the following kernel warning: | | | |
| | | | | | |
| | | tracepoint_probe_unregister+0x280/0x370 [] | | | |
| | | CPU: 0 PID: 1217 Comm: trace-cmd Not tainted 5.17.0-rc6-next-20220307-nico+ #19 | | | |
| | | RIP: 0010:tracepoint_probe_unregister+0x280/0x370 [] | | | |
| | | CR2: 00007ff919b29497 CR3: 0000000109da4005 CR4: 000000000170ef0 | | | |
| | | Call Trace: <task></task> | | | |
| | | osnoise_workload_stop+0x36/0x90 | | | |

| | | tracing_set_tracer+0x108/0x260 tracing_set_trace_write+0x94/0xd0 | | | |
|----------------|-------------------------------|--|------------|-----|------|
| | | ?check_object_size.part.0+0x10a/0x150 ? selinux_file_permission+0x104/0x150 vfs_write+0xb5/0x290 ksys_write+0x5f/0xe0 do_syscall_64+0x3b/0x90 | | | |
| | | entry_SYSCALL_64_after_hwframe+0x44/0xae RIP: 0033:0x7ff919a18127 [] | | | |
| | | [end trace 000000000000000000000000000000000000 | | | |
| | | The warning complains about an attempt to unregister an unregistered tracepoint. | | | |
| | | This happens on trace-cmd because it first stops tracing, and then switches the tracer to nop. Which is equivalent to: | | | |
| | | # cd /sys/kernel/tracing/ # echo osnoise > current_tracer # echo 0 > tracing_on # echo nop > current_tracer | | | |
| | | The osnoise tracer stops the workload when no trace instance is actually collecting data. This can be caused both by disabling tracing or disabling the tracer itself. | | | |
| | | To avoid unregistering events twice, use the existing trace_osnoise_callback_enabled variable to check if the events (and the workload) are actually active before trying to deactivate them. | | | |
| CVE-2022-48851 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.8 | High |
| | | staging: gdm724x: fix use after free in gdm_lte_rx() | | | |
| | | The netif_rx_ni() function frees the skb so we can't dereference it to save the skb->len. | | | |
| CVE-2022-48854 | linux - linux_kernel | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.8 | High |
| | | net: arc_emac: Fix use after free in arc_mdio_probe() | | | |
| | | If bus->state is equal to MDIOBUS_ALLOCATED, mdiobus_free(bus) will free the "bus". But bus->name is still used in the next line, which will lead | | | |
| | | to a use after free. | | | |
| | | We can fix it by putting the name in a local variable and make the bus->name point to the rodata section "name",then use the name in the | | | |
| CVE-2024-39731 | ibm - multiple | error message without referring to bus to avoid the uaf. IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 uses | 2024-07-15 | 7.5 | High |
| | products | weaker than expected cryptographic algorithms that could allow an attacker to decrypt highly sensitive information. IBM X-Force ID: 295970. | | | J |
| CVE-2024-21175 | oracle - multiple products | Vulnerability in the Oracle WebLogic Server product of Oracle Fusion Middleware (component: Core). Supported versions that are affected are 12.2.1.4.0 and 14.1.1.0.0. Easily exploitable vulnerability allows unauthenticated attacker with network access via HTTP to compromise Oracle WebLogic Server. Successful | 2024-07-16 | 7.5 | High |
| | | attacks of this vulnerability can result in unauthorized creation, deletion or modification access to critical data or all Oracle WebLogic Server accessible data. CVSS 3.1 Base Score 7.5 (Integrity impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:N/UI:N/S:U/C:N/I:H/A:N). | | | |
| CVE-2024-21182 | oracle - multiple products | Vulnerability in the Oracle WebLogic Server product of Oracle Fusion Middleware (component: Core). Supported versions that are affected are 12.2.1.4.0 and 14.1.1.0.0. Easily exploitable | 2024-07-16 | 7.5 | High |
| | | vulnerability allows unauthenticated attacker with network access via T3, IIOP to compromise Oracle WebLogic Server. Successful attacks of this vulnerability can result in unauthorized access to critical data or complete access to all Oracle WebLogic Server accessible data. CVSS 3.1 Base Score 7.5 (Confidentiality impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:N/A:N). | | | |
| CVE-2024-21183 | oracle - multiple products | Vulnerability in the Oracle WebLogic Server product of Oracle Fusion Middleware (component: Core). Supported versions that are affected are 12.2.1.4.0 and 14.1.1.0.0. Easily exploitable vulnerability allows unauthenticated attacker with network access | 2024-07-16 | 7.5 | High |
| | | via T3, IIOP to compromise Oracle WebLogic Server. Successful | | | |

| | | attacks of this vulnerability can result in unauthorized access to critical data or complete access to all Oracle WebLogic Server accessible data. CVSS 3.1 Base Score 7.5 (Confidentiality impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:N/A:N). | | | |
|----------------|-------------------------------|---|------------|-----|------|
| CVE-2024-32007 | apache - multiple products | An improper input validation of the p2c parameter in the Apache CXF JOSE code before 4.0.5, 3.6.4 and 3.5.9 allows an attacker to perform a denial of service attack by specifying a large value for this parameter in a token. | 2024-07-19 | 7.5 | High |
| CVE-2024-21184 | oracle - database_server | Vulnerability in the Oracle Database RDBMS Security component of Oracle Database Server. Supported versions that are affected are 19.3-19.23. Easily exploitable vulnerability allows high privileged attacker having Execute on SYS.XS_DIAG privilege with network access via Oracle Net to compromise Oracle Database RDBMS Security. Successful attacks of this vulnerability can result in takeover of Oracle Database RDBMS Security. CVSS 3.1 Base Score 7.2 (Confidentiality, Integrity and Availability impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:H/UI:N/S:U/C:H/I:H/A:H). | 2024-07-16 | 7.2 | High |
| CVE-2022-48855 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.1 | High |
| | | sctp: fix kernel-infoleak for SCTP sockets | | | |
| | | syzbot reported a kernel infoleak [1] of 4 bytes. | | | |
| | | After analysis, it turned out r->idiag_expires is not initialized if inet_sctp_diag_fill() calls inet_diag_msg_common_fill() | | | |
| | | Make sure to clear idiag_timer/idiag_retrans/idiag_expires and let inet_diag_msg_sctpasoc_fill() fill them again if needed. | | | |
| | | [1] | | | |
| | | BUG: KMSAN: kernel-infoleak in instrument_copy_to_user include/linux/instrumented.h:121 [inline] BUG: KMSAN: kernel-infoleak in copyout lib/iov_iter.c:154 [inline] BUG: KMSAN: kernel-infoleak in _copy_to_iter+0x6ef/0x25a0 lib/iov_iter.c:668 instrument_copy_to_user include/linux/instrumented.h:121 [inline] copyout lib/iov_iter.c:154 [inline] _copy_to_iter+0x6ef/0x25a0 lib/iov_iter.c:668 copy_to_iter include/linux/uio.h:162 [inline] simple_copy_to_iter+0xf3/0x140 net/core/datagram.c:519skb_datagram_iter+0x2d5/0x11b0 net/core/datagram.c:425 skb_copy_datagram_iter+0xdc/0x270 net/core/datagram.c:533 | | | |
| | | skb_copy_datagram_msg include/linux/skbuff.h:3696 [inline] netlink_recvmsg+0x669/0x1c80 net/netlink/af_netlink.c:1977 sock_recvmsg_nosec net/socket.c:948 [inline] sock_recvmsg net/socket.c:966 [inline]sys_recvfrom+0x795/0xa10 net/socket.c:2097do_sys_recvfrom net/socket.c:2115 [inline]se_sys_recvfrom net/socket.c:2111 [inline]x64_sys_recvfrom+0x19d/0x210 net/socket.c:2111 do_syscall_x64 arch/x86/entry/common.c:51 [inline] do_syscall_64+0x54/0xd0 arch/x86/entry/common.c:82 entry_SYSCALL_64_after_hwframe+0x44/0xae | | | |
| | | Uninit was created at: slab_post_alloc_hook mm/slab.h:737 [inline] slab_alloc_node mm/slub.c:3247 [inline]kmalloc_node_track_caller+0xe0c/0x1510 mm/slub.c:4975 kmalloc_reserve net/core/skbuff.c:354 [inline]alloc_skb+0x545/0xf90 net/core/skbuff.c:426 alloc_skb include/linux/skbuff.h:1158 [inline] netlink_dump+0x3e5/0x16c0 net/netlink/af_netlink.c:2248netlink_dump_start+0xcf8/0xe90 net/netlink/af_netlink.c:2373 netlink_dump_start include/linux/netlink.h:254 [inline] inet_diag_handler_cmd+0x2e7/0x400 net/ipv4/inet_diag.c:1341 sock_diag_rcv_msg+0x24a/0x620 netlink_rcv_skb+0x40c/0x7e0 net/netlink/af_netlink.c:2494 sock_diag_rcv+0x63/0x80 net/core/sock_diag.c:277 netlink_unicast_kernel net/netlink/af_netlink.c:1317 [inline] | | | |
| | | netlink_unicast+0x1093/0x1360 net/netlink/af_netlink.c:1343 netlink_sendmsg+0x14d9/0x1720 net/netlink/af_netlink.c:1919 sock_sendmsg_nosec net/socket.c:705 [inline] sock_sendmsg net/socket.c:725 [inline] sock_write_iter+0x594/0x690 net/socket.c:1061 do_iter_readv_writev+0xa7f/0xc70 do_iter_write+0x52c/0x1500 fs/read_write.c:851 vfs_writev fs/read_write.c:924 [inline] | | | |

| | | do_writev+0x645/0xe00 fs/read_write.c:967 | | | |
|-----------------------|------------------------------|---|------------|-----|---------|
| | | do_sys_writev fs/read_write.c:1040 [inline] | | | |
| | | se_sys_writev fs/read_write.c:1037 [inline] | | | |
| | | x64_sys_writev+0xe5/0x120 fs/read_write.c:1037 | | | |
| | | do_syscall_x64 arch/x86/entry/common.c:51 [inline] | | | |
| | | do_syscall_64+0x54/0xd0 arch/x86/entry/common.c:82 entry SYSCALL 64 after hwframe+0x44/0xae | | | |
| | | entry_SYSCALL_64_arter_nwrrame+0x44/0xae | | | |
| | | Bytes 68-71 of 2508 are uninitialized | | | |
| | | Memory access of size 2508 starts at ffff888114f9b000 | | | |
| | | Data copied to user address 00007f7fe09ff2e0 | | | |
| | | | | | |
| | | CPU: 1 PID: 3478 Comm: syz-executor306 Not tainted 5.17.0-rc4- | | | |
| | | syzkaller #0 | | | |
| | | Hardware name: Google Google Compute Engine/Google Compute | | | |
| CV /F 2022 40066 | P 100 d. | Engine, BIOS Google 01/01/2011 | 2024.07.46 | 7.4 | 112.1. |
| CVE-2022-48866 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7.1 | High |
| | products | HID: hid-thrustmaster: fix OOB read in thrustmaster_interrupts | | | |
| | | This. The thrustinaster. his oob read in thrustinaster_interrupts | | | |
| | | Syzbot reported an slab-out-of-bounds Read in | | | |
| | | thrustmaster_probe() bug. | | | |
| | | The root case is in missing validation check of actual number of | | | |
| | | endpoints. | | | |
| | | | | | |
| | | Code should not blindly access usb_host_interface::endpoint | | | |
| | | array, since | | | |
| | | it may contain less endpoints than code expects. | | | |
| | | Fix it by adding missing validation about and print an array if | | | |
| | | Fix it by adding missing validaion check and print an error if number of endpoints do not match expected number | | | |
| CVE-2022-48858 | linux - multiple | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 7 | High |
| <u>CVE 2022 10030</u> | products | The Emax Reffici, the following value ability has been resolved. | 20210710 | , | 1 11811 |
| | F | net/mlx5: Fix a race on command flush flow | | | |
| | | | | | |
| | | Fix a refcount use after free warning due to a race on command | | | |
| | | entry. | | | |
| | | Such race occurs when one of the commands releases its last | | | |
| | | refcount and | | | |
| | | frees its index and entry while another process running command | | | |
| | | flush | | | |
| | | flow takes refcount to this command entry. The process which handles | | | |
| | | commands flush may see this command as needed to be flushed if | | | |
| | | the other | | | |
| | | process released its refcount but didn't release the index yet. Fix it | | | |
| | | by adding the needed spin lock. | | | |
| | | | | | |
| | | It fixes the following warning trace: | | | |
| | | | | | |
| | | refcount_t: addition on 0; use-after-free. | | | |
| | | WARNING: CPU: 11 PID: 540311 at lib/refcount.c:25 | | | |
| | | refcount_warn_saturate+0x80/0xe0 | | | |
| | | RIP: 0010:refcount_warn_saturate+0x80/0xe0 | | | |
| | | | | | |
| | | Call Trace: | | | |
| | | <task></task> | | | |
| | | mlx5_cmd_trigger_completions+0x293/0x340 [mlx5_core] | | | |
| | | mlx5_cmd_flush+0x3a/0xf0 [mlx5_core] | | | |
| | | enter_error_state+0x44/0x80 [mlx5_core] | | | |
| | | mlx5_fw_fatal_reporter_err_work+0x37/0xe0 [mlx5_core] | | | |
| | | process_one_work+0x1be/0x390 | | | |
| | | worker_thread+0x4d/0x3d0 | | | |
| | | ? rescuer_thread+0x350/0x350 kthread+0x141/0x160 | | | |
| | | ? set_kthread_struct+0x40/0x40 | | | |
| | | ret_from_fork+0x1f/0x30 | | | |
| | | | | | |
| CVE-2023-41916 | apache - linkis | | 2024-07-15 | 6.5 | Medium |
| | | In Apache Linkis =1.4.0, due to the lack of effective filtering | | | |
| | | of parameters, an attacker configuring malicious Mysql JDBC | | | |
| | | parameters in the DataSource Manager Module will | | | |
| | | trigger arbitrary file reading. Therefore, the parameters in the | | | |
| | | Mysql JDBC URL should be blacklisted. This attack requires the attacker to obtain an authorized account from Linkis before it can | | | |
| | | be carried out. Versions of Apache Linkis = 1.4.0 will be affected. | | | |
| | | We recommend users upgrade the version of Linkis to version | | | |
| | | 1.5.0. | | | |
| | | | | | |
| | | | | | |

| CVE-2024-21177 | oracle - multiple | Vulnerability in the MySQL Server product of Oracle MySQL | 2024-07-16 | 6.5 | Medium |
|-----------------------|-------------------|--|------------|-----|-----------|
| | products | (component: Server: Optimizer). Supported versions that are | | | |
| | | affected are 8.0.37 and prior and 8.4.0 and prior. Easily | | | |
| | | exploitable vulnerability allows low privileged attacker with network access via multiple protocols to compromise MySQL | | | |
| | | Server. Successful attacks of this vulnerability can result in | | | |
| | | unauthorized ability to cause a hang or frequently repeatable | | | |
| | | crash (complete DOS) of MySQL Server. CVSS 3.1 Base Score 6.5 | | | |
| | | (Availability impacts). CVSS Vector: | | | |
| CVE 2024 2004 | | (CVSS:3.1/AV:N/AC:L/PR:L/UI:N/S:U/C:N/I:N/A:H). | 2024.07.16 | 6.5 | N A = ali |
| CVE-2024-2884 | google - chrome | Out of bounds read in V8 in Google Chrome prior to 121.0.6167.139 allowed a remote attacker to potentially perform | 2024-07-16 | 6.5 | Medium |
| | | out of bounds memory access via a crafted HTML page. | | | |
| | | (Chromium security severity: Medium) | | | |
| CVE-2024-5500 | google - chrome | Inappropriate implementation in Sign-In in Google Chrome prior to | 2024-07-16 | 6.5 | Medium |
| | | 1.3.36.351 allowed a remote attacker to bypass navigation | | | |
| | | restrictions via a crafted HTML page. (Chromium security severity: Medium) | | | |
| CVE-2024-3175 | google - chrome | Insufficient data validation in Extensions in Google Chrome prior to | 2024-07-16 | 6.3 | Medium |
| <u>CVE 202 3173</u> | google emonie | 120.0.6099.62 allowed a remote attacker to perform privilege | 20210710 | 0.5 | Wicaian |
| | | escalation via a crafted Chrome Extension. (Chromium security | | | |
| | | severity: Low) | | | |
| CVE-2024-21178 | oracle - multiple | Vulnerability in the PeopleSoft Enterprise PeopleTools product of | 2024-07-16 | 6.1 | Medium |
| | products | Oracle PeopleSoft (component: Portal). Supported versions that | | | |
| | | are affected are 8.59, 8.60 and 8.61. Easily exploitable | | | |
| | | vulnerability allows unauthenticated attacker with network access via HTTP to compromise PeopleSoft Enterprise PeopleTools. | | | |
| | | Successful attacks require human interaction from a person other | | | |
| | | than the attacker and while the vulnerability is in PeopleSoft | | | |
| | | Enterprise PeopleTools, attacks may significantly impact additional | | | |
| | | products (scope change). Successful attacks of this vulnerability | | | |
| | | can result in unauthorized update, insert or delete access to some | | | |
| | | of PeopleSoft Enterprise PeopleTools accessible data as well as unauthorized read access to a subset of PeopleSoft Enterprise | | | |
| | | PeopleTools accessible data. CVSS 3.1 Base Score 6.1 | | | |
| | | (Confidentiality and Integrity impacts). CVSS Vector: | | | |
| | | (CVSS:3.1/AV:N/AC:L/PR:N/UI:R/S:C/C:L/I:L/A:N). | | | |
| CVE-2024-21188 | oracle - multiple | Vulnerability in the Oracle Financial Services Revenue | 2024-07-16 | 6.1 | Mediun |
| | products | Management and Billing product of Oracle Financial Services | | | |
| | | Applications (component: Chatbot). Supported versions that are | | | |
| | | affected are 6.0.0.0.0 and 6.1.0.0.0. Easily exploitable vulnerability allows unauthenticated attacker with network access via HTTP to | | | |
| | | compromise Oracle Financial Services Revenue Management and | | | |
| | | Billing. Successful attacks require human interaction from a | | | |
| | | person other than the attacker and while the vulnerability is in | | | |
| | | Oracle Financial Services Revenue Management and Billing, attacks | | | |
| | | may significantly impact additional products (scope change). | | | |
| | | Successful attacks of this vulnerability can result in unauthorized update, insert or delete access to some of Oracle Financial Services | | | |
| | | Revenue Management and Billing accessible data as well as | | | |
| | | unauthorized read access to a subset of Oracle Financial Services | | | |
| | | Revenue Management and Billing accessible data. CVSS 3.1 Base | | | |
| | | Score 6.1 (Confidentiality and Integrity impacts). CVSS Vector: | | | |
| CVE 2022 49925 | linux - multiple | (CVSS:3.1/AV:N/AC:L/PR:N/UI:R/S:C/C:L/I:L/A:N). | 2024-07-16 | 5.5 | Medium |
| CVE-2022-48835 | products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | iviedium |
| | products | scsi: mpt3sas: Page fault in reply q processing | | | |
| | | | | | |
| | | A page fault was encountered in mpt3sas on a LUN reset error | | | |
| | | path: | | | |
| | | [145 762216] mnt2cas, cm1: Task abort to failed: | | | |
| | | [145.763216] mpt3sas_cm1: Task abort tm failed: handle(0x0002),timeout(30) tr_method(0x0) smid(3) | | | |
| | | msix_index(0) | | | |
| | | [145.778932] scsi 1:0:0:0: task abort: FAILED | | | |
| | | scmd(0x000000024ba29a2) | | | |
| | | [145.817307] scsi 1:0:0:0: attempting device reset! | | | |
| | | scmd(0x000000024ba29a2) | | | |
| | | [145.827253] scsi 1:0:0:0: [sg1] tag#2 CDB: Receive Diagnostic 1c 01 01 ff fc 00 | | | |
| | | [145.837617] scsi target1:0:0: handle(0x0002), | | | |
| | | sas_address(0x500605b0000272b9), phy(0) | | | |
| | | [145.848598] scsi target1:0:0: enclosure logical | | | |
| | | id(0x500605b0000272b8), slot(0) | | | |
| | | [149.858378] mpt3sas_cm1: Poll ReplyDescriptor queues for | | | |
| | | completion of smid(0), task_type(0x05), handle(0x0002) | | | |
| | | | | | |
| | | [149.875202] BUG: unable to handle page fault for address: 00000007fffc445d | | | |

| [149,89313] PGD 0 P4D 0 [149,90312] PGD 0 P4D 0 [149,90313] PGD 0 P4D 0 [149,90313] PGD 0:000 [H1] PREEMPT SMP NOPTI [149,917417] CPU: 24 PID: 3512 Comm: scsi_eh_1 Kdump: loaded Tainted: 65 0 5.10,89-attav-1 H1 [149,917417] CPU: 24 PID: 3512 Comm: scsi_eh_1 Kdump: loaded Tainted: 65 0 5.10,89-attav-1 H1 [149,934327] Hardware name: DDN 200NVX2 [7200NVX2-MB | |
|--|--------|
| CR4: 0000000000770ee0 [150.086887] DR0: 000000000000000 DR1: 000000000000000 DR2: 0000000000000000 [150.097670] DR3: 0000000000000000 DR6: 00000000fffe0ff0 DR7: 0000000000000000 [150.108323] PKRU: 55555554 [150.114690] Call Trace: [150.120497] ? printk+0x48/0x4a [150.127049] mpt3sas_scsih_issue_tm.cold.114+0x2e/0x2b3 [mpt3sas] [150.136453] mpt3sas_scsih_issue_locked_tm+0x86/0xb0 [mpt3sas] [150.13645759] scsih_dev_reset+0xea/0x300 [mpt3sas] [150.153891] scsi_eh_ready_devs+0x541/0x9e0 [scsi_mod] [150.162206] ? _scsi_host_match+0x20/0x20 [scsi_mod] [150.178925] ? blk_mq_tagset_busy_iter+0x45/0x60 [150.186638] ? scsi_try_target_reset+0x90/0x90 [scsi_mod] [150.195087] scsi_etnor_handler+0x3a5/0x4a0 [scsi_mod] [150.203206] ? _schedule+0x1e9/0x610 [150.203206] ? _schedule+0x1e9/0x510 [150.217924] kthread+0x12e/0x150 [150.224041] ? kthread+0x12e/0x150 [150.231206] ret_from_fork+0x1f/0x30 This is caused by mpt3sas_base_sync_reply_irqs() using an invalid reply_q pointer outside of the list_for_each_entry() loop. At the end of the full list traversal the pointer is invalid. | |
| Move the _base_process_reply_queue() call inside of the loop. ultiple | Лedium |
| endpoint type. There was a check for the number of endpoints, but not for the type of endpoint. Fix it by replacing old desc.bNumEndpoints check with usb_find_common_endpoints() helper for finding endpoints Fail log: usb 5-1: BOGUS urb xfer, pipe 1 != type 3 | |
| WARNING: CPU: 2 PID: 48 at drivers/usb/core/urb.c:502 usb_submit_urb+0xed2/0x18a0 drivers/usb/core/urb.c:502 Modules linked in: | |

| | | T | Т | I | 1 |
|----------------|------------------|---|------------|-----|-----------|
| | | CPU: 2 PID: 48 Comm: kworker/2:2 Not tainted 5.17.0-rc6-syzkaller-00226-g07ebd38a0da2 #0 Hardware name: QEMU Standard PC (Q35 + ICH9, 2009), BIOS 1.14.0-2 04/01/2014 | | | |
| | | Workqueue: usb_hub_wq hub_event | | | |
| | | Call Trace: | | | |
| | | <task> aiptek_open+0xd5/0x130 drivers/input/tablet/aiptek.c:830 input_open_device+0x1bb/0x320 drivers/input/input.c:629</task> | | | |
| CVE-2022-48838 | linux - multiple | kbd_connect+0xfe/0x160 drivers/tty/vt/keyboard.c:1593 In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | usb: gadget: Fix use-after-free bug by not setting udc->dev.driver | | | |
| | | The syzbot fuzzer found a use-after-free bug: | | | |
| | | BUG: KASAN: use-after-free in dev_uevent+0x712/0x780 drivers/base/core.c:2320 Read of size 8 at addr ffff88802b934098 by task udevd/3689 | | | |
| | | CPU: 2 PID: 3689 Comm: udevd Not tainted 5.17.0-rc4-syzkaller-00229-g4f12b742eb2b #0 Hardware name: QEMU Standard PC (Q35 + ICH9, 2009), BIOS 1.14.0-2 04/01/2014 | | | |
| | | Call Trace: <task> dump_stack lib/dump_stack.c:88 [inline] dump_stack_lvl+0xcd/0x134 lib/dump_stack.c:106 print_address_description.constprop.0.cold+0x8d/0x303 mm/kasan/report.c:255 kasan_report mm/kasan/report.c:442 [inline] kasan_report.cold+0x83/0xdf mm/kasan/report.c:459 dev_uevent+0x712/0x780 drivers/base/core.c:2320 uevent_show+0x1b8/0x380 drivers/base/core.c:2391</task> | | | |
| | | dev_attr_show+0x4b/0x90 drivers/base/core.c:2094 | | | |
| | | Although the bug manifested in the driver core, the real cause was a race with the gadget core. dev uevent() does: | | | |
| | | <pre>if (dev->driver)</pre> | | | |
| | | and between the test and the dereference of dev->driver, the gadget core sets dev->driver to NULL. | | | |
| | | The race wouldn't occur if the gadget core registered its devices on a real bus, using the standard synchronization techniques of the driver core. However, it's not necessary to make such a large change in order to fix this bug; all we need to do is make sure that udc->dev.driver is always NULL. | | | |
| | | In fact, there is no reason for udc->dev.driver ever to be set to anything, let alone to the value it currently gets: the address of the gadget's driver. After all, a gadget driver only knows how to manage a gadget, not how to manage a UDC. | | | |
| CVE-2022-48839 | linux - multiple | This patch simply removes the statements in the gadget core that touch udc->dev.driver. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| <u> </u> | products | net/packet: fix slab-out-of-bounds access in packet_recvmsg() | 2024 07-10 | 5.5 | Wiculuiii |
| | | syzbot found that when an AF_PACKET socket is using PACKET_COPY_THRESH and mmap operations, tpacket_rcv() is queueing skbs with | | | |
| | | garbage in skb->cb[], triggering a too big copy [1] Presumably, users of af_packet using mmap() already gets correct metadata from the mapped buffer, we can simply make sure to clear 12 bytes that might be copied to user space later. | | | |
| | | BUG: KASAN: stack-out-of-bounds in memcpy include/linux/fortify-string.h:225 [inline] | | | |

| | Tana many a second | T | | |
|--------------------------------------|--|------------|-----|--------|
| | BUG: KASAN: stack-out-of-bounds in packet_recvmsg+0x56c/0x1150 net/packet/af_packet.c:3489 Write of size 165 at addr ffffc9000385fb78 by task syzexecutor233/3631 | | | |
| | CPU: 0 PID: 3631 Comm: syz-executor233 Not tainted 5.17.0-rc7-syzkaller-02396-g0b3660695e80 #0 | | | |
| | Hardware name: Google Google Compute Engine/Google Compute Engine, BIOS Google 01/01/2011 Call Trace: | | | |
| | <task>dump_stack lib/dump_stack.c:88 [inline] dump_stack_lvl+0xcd/0x134 lib/dump_stack.c:106</task> | | | |
| | print_address_description.constprop.0.cold+0xf/0x336 mm/kasan/report.c:255kasan_report mm/kasan/report.c:442 [inline] kasan_report.cold+0x83/0xdf mm/kasan/report.c:459 | | | |
| | check_region_inline mm/kasan/generic.c:183 [inline] kasan_check_range+0x13d/0x180 mm/kasan/generic.c:189 memcpy+0x39/0x60 mm/kasan/shadow.c:66 | | | |
| | memcpy include/linux/fortify-string.h:225 [inline] packet_recvmsg+0x56c/0x1150 net/packet/af_packet.c:3489 sock_recvmsg_nosec net/socket.c:948 [inline] | | | |
| | sock_recvmsg net/socket.c:966 [inline] sock_recvmsg net/socket.c:962 [inline]sys_recvmsg+0x2c4/0x600 net/socket.c:2632sys_recvmsg+0x127/0x200 net/socket.c:2674 | | | |
| | sys_recvmsg+0xe2/0x1a0 net/socket.c:2704 do_syscall_x64 arch/x86/entry/common.c:50 [inline] do_syscall_64+0x35/0xb0 arch/x86/entry/common.c:80 | | | |
| | entry_SYSCALL_64_after_hwframe+0x44/0xae RIP: 0033:0x7fdfd5954c29 Code: 28 00 00 00 75 05 48 83 c4 28 c3 e8 41 15 00 00 90 48 89 f8 48 89 f7 48 89 d6 48 89 ca 4d 89 c2 4d 89 c8 4c 8b 4c 24 08 0f 05 | | | |
| | 48 89 17 48 89 00 48 89 03 40 69 02 40 89 08 40 80 40 03 48 89 01 03 48 89 01 01 03 48 89 01 00 00000000000000000000000000000 | | | |
| | RAX: fffffffffffda RBX: 00000000000000 RCX: 00007fdfd5954c29 RDX: 00000000000000 RSI: 000000020000500 RDI: | | | |
| | 000000000000005 RBP: 00000000000000 R08: 0000000000000 R09: 0000000000000000 R10: 00000000000000 R11: 000000000000246 R12: 00007ffcf8e71e60 | | | |
| | R13: 0000000000f4240 R14: 000000000000c1ff R15: 00007ffcf8e71e54 | | | |
| | addr ffffc9000385fb78 is located in stack of task syz- executor233/3631 at offset 32 in frame: sys_recvmsg+0x0/0x600 include/linux/uio.h:246 | | | |
| | this frame has 1 object: [32, 160) 'addr' | | | |
| | Memory state around the buggy address: ffffc9000385fa80: 00 04 f3 f3 f3 f3 00 00 00 00 00 00 00 00 ffffc9000385fb00: 00 00 00 00 00 00 00 00 00 00 f1 f1 f1 f1 f1 00 >ffffc9000385fb80: 00 00 00 00 00 00 00 00 00 00 00 00 0 | | | |
| | ffffc9000385fc00: f3 f3 f3 00 00 00 00 00 00 00 00 00 00 00 f1 ffffc9000385fc80: f1 f1 f1 00 f2 f2 f2 00 f2 f2 f2 00 00 00 00 00 ========== | | | |
| CVE-2022-48840 linux - mult products | , | 2024-07-16 | 5.5 | Medium |
| | iavf: Fix hang during reboot/shutdown | | | |
| | Recent commit 974578017fc1 ("iavf: Add waiting so the port is initialized in remove") adds a wait-loop at the beginning of iavf_remove() to ensure that port initialization is finished prior unregistering net device. This causes a regression | | | |
| | in reboot/shutdown scenario because in this case callback iavf_shutdown() is called and this callback detaches the device, makes it down if it is running and sets its state toIAVF_REMOVE. | | | |
| | Later shutdown callback of associated PF driver (e.g. ice_shutdown) | | | |

| | | | | 1 | , |
|----------------|------------------------------|--|------------|-----|--------|
| | | is called. That callback calls among other things sriov_disable() that calls indirectly iavf_remove() (see stack trace below). As the adapter state is alreadyIAVF_REMOVE then the | | | |
| | | mentioned loop is end-less and shutdown process hangs. | | | |
| | | The patch fixes this by checking adapter's state at the beginning of iavf_remove() and skips the rest of the function if the adapter is already in remove state (shutdown is in progress). | | | |
| | | Reproducer: 1. Create VF on PF driven by ice or i40e driver 2. Ensure that the VF is bound to iavf driver 3. Reboot | | | |
| | | [52625.981294] sysrq: SysRq : Show Blocked State [52625.988377] task:reboot state:D stack: 0 pid:17359 ppid: 1 f2 | | | |
| | | [52625.996732] Call Trace: [52625.999187]schedule+0x2d1/0x830 [52626.007400] schedule+0x35/0xa0 [52626.010545] schedule_hrtimeout_range_clock+0x83/0x100 | | | |
| | | [52626.020046] usleep_range+0x5b/0x80 [52626.023540] iavf_remove+0x63/0x5b0 [iavf] [52626.027645] pci_device_remove+0x3b/0xc0 | | | |
| | | [52626.031572] device_release_driver_internal+0x103/0x1f0 [52626.036805] pci_stop_bus_device+0x72/0xa0 [52626.040904] pci_stop_and_remove_bus_device+0xe/0x20 | | | |
| | | [52626.045870] pci_iov_remove_virtfn+0xba/0x120 [52626.050232] sriov_disable+0x2f/0xe0 [52626.053813] ice_free_vfs+0x7c/0x340 [ice] [52626.057946] ice_remove+0x220/0x240 [ice] | | | |
| | | [52626.061967] ice_shutdown+0x16/0x50 [ice] [52626.065987] pci_device_shutdown+0x34/0x60 [52626.070086] device_shutdown+0x165/0x1c5 | | | |
| | | [52626.074011] kernel_restart+0xe/0x30 [52626.077593]do_sys_reboot+0x1d2/0x210 [52626.093815] do_syscall_64+0x5b/0x1a0 | | | |
| CVE-2022-48841 | linux - multiple | [52626.097483] entry_SYSCALL_64_after_hwframe+0x65/0xca In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | ice: fix NULL pointer dereference in ice_update_vsi_tx_ring_stats() | | | |
| | | It is possible to do NULL pointer dereference in routine that | | | |
| | | updates Tx ring stats. Currently only stats and bytes are updated when ring pointer is valid, but later on ring is accessed to propagate gathered | | | |
| | | Tx stats onto VSI stats. | | | |
| CVE-2022-48843 | linux - multiple products | Change the existing logic to move to next ring when ring is NULL. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | • | drm/vrr: Set VRR capable prop only if it is attached to connector | | | |
| | | VRR capable property is not attached by default to the connector It is attached only if VRR is supported. So if the driver tries to call drm core set prop function without it being attached that causes NULL dereference. | | | |
| CVE-2022-48844 | linux - multiple | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | Bluetooth: hci_core: Fix leaking sent_cmd skb | | | |
| | | sent_cmd memory is not freed before freeing hci_dev causing it to leak it contents. | | | |
| CVE-2022-48845 | linux - multiple | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | MIPS: smp: fill in sibling and core maps earlier | | | |
| | | After enabling CONFIG_SCHED_CORE (landed during 5.14 cycle), 2-core 2-thread-per-core interAptiv (CPS-driven) started emitting the following: | | | |
| | | [0.025698] CPU1 revision is: 0001a120 (MIPS interAptiv (multi)) [0.048183][cut here] [0.048187] WARNING: CPU: 1 PID: 0 at | | | |
| | | kernel/sched/core.c:6025 sched_core_cpu_starting+0x198/0x240 [0.048220] Modules linked in: | | | |

| performing delay calibration, but the WARN splat is now gone. The very same debug prints now yield exactly what I expected from them: [0.048433] CPU: 1, smt_mask: 0-1 [0] https://git.kernel.org/pub/scm/linux/kernel/git/mips/linux.git/co mmit/?id=76ce7cfe35ef | | | [0.048233] CPU: 1 PID: 0 Comm: swapper/1 Not tainted 5.17.0-rc3+#35 b7b319f24073fd9a3c2aa7ad15fb7993eec0b26f [0.048247] Stack : 817f0000 00000004 327804c8 810eb050 00000000 000000000 00000000000000 | | | |
|--|----------------|-----|---|------------|-----|--------|
| [0] https://git.kernel.org/pub/scm/linux/kernel/git/mips/linux.git/co | | | While the flow is same on ARM64 (maps after the notifier, although before calling set_cpu_online()), x86 started calculating sibling maps earlier than starting the CPU HP callbacks in Linux 4.14 (see [0] for the reference). Neither me nor my brief tests couldn't find any potential caveats in calculating the maps right after performing delay calibration, but the WARN splat is now gone. The very same debug prints now yield exactly what I expected from | | | |
| 4 | | | [0] | | | |
| CVE-2022-48846 linux - multiple products lin | CVE-2022-48846 | · · | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| block: release rq qos structures for queue without disk | | , | block: release rq qos structures for queue without disk | | | |
| blkcg_init_queue() may add rq qos structures to request queue, | | | | | | |
| previously blk_cleanup_queue() calls rq_qos_exit() to release them, but | | | blk_cleanup_queue() calls rq_qos_exit() to release them, but | | | |
| commit 8e141f9eb803 ("block: drain file system I/O on del_gendisk") | | | 8e141f9eb803 ("block: drain file system I/O on del_gendisk") | | | |
| moves rq_qos_exit() into del_gendisk(), so memory leak is caused because queues may not have disk, such as un-present scsi luns, | | | ·-· | | | |
| nvme admin queue, | | | nvme | | | |

| products drm/amdgpu: bypass tiling flag check in virtual display case (v2) vkms leverages common amdgpu framebuffer creation, and also as it does not support FB modifier, there is no need to check tiling flags when initing framebuffer when virtual display is enabled. This can fix below calltrace: amdgpu 0000:00:08.0: GFX9+ requires FB check based on format modifier WARNING: CPU: 0 PID: 1023 at drivers/gpu/drm/amd/amdgpu/amdgpu_display.c:1150 amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. | edium |
|---|-------|
| blkcg_init_queue()/blkcg_exit_queue() into disk allocation/release handler, and patches have been in for-5.18/block. CVE-2022-48849 linux - multiple products blkcg_init_queue()/blkcg_exit_queue() into disk allocation/release handler, and patches have been in for-5.18/block. In the Linux kernel, the following vulnerability has been resolved: linux - multiple products linux - multiple products linux - multiple products blkcg_init_queue()/blkcg_exit_queue() into disk allocation/release handler, and patches have been in for-5.18/block. In the Linux kernel, the following vulnerability has been resolved: linux - multiple products linux - multiple products | |
| In the Linux kernel, the following vulnerability has been resolved: 2024-07-16 5.5 M | |
| drm/amdgpu: bypass tiling flag check in virtual display case (v2) vkms leverages common amdgpu framebuffer creation, and also as it does not support FB modifier, there is no need to check tiling flags when initing framebuffer when virtual display is enabled. This can fix below calltrace: amdgpu 0000:00:08.0: GFX9+ requires FB check based on format modifier WARNING: CPU: 0 PID: 1023 at drivers/gpu/drm/amd/amdgpu/amdgpu_display.c:1150 amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. CVE-2022-48850 linux - multiple products linux - multiple products In the Linux kernel, the following vulnerability has been resolved: 0.55 N net-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| also as it does not support FB modifier, there is no need to check tiling flags when initing framebuffer when virtual display is enabled. This can fix below calltrace: amdgpu 0000:00:08.0: GFX9+ requires FB check based on format modifier WARNING: CPU: 0 PID: 1023 at drivers/gpu/drm/amd/amdgpu/amdgpu_display.c:1150 amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. CVE-2022-48850 linux - multiple products In the Linux kernel, the following vulnerability has been resolved: 2024-07-16 5.5 Ment-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| amdgpu 0000:00:08.0: GFX9+ requires FB check based on format modifier WARNING: CPU: 0 PID: 1023 at drivers/gpu/drm/amd/amdgpu/amdgpu_display.c:1150 amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. CVE-2022-48850 Iinux - multiple products In the Linux kernel, the following vulnerability has been resolved: 2024-07-16 5.5 M When bringing down the netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| modifier WARNING: CPU: 0 PID: 1023 at drivers/gpu/drm/amd/amdgpu_display.c:1150 amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. CVE-2022-48850 linux - multiple products In the Linux kernel, the following vulnerability has been resolved: net-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| amdgpu_display_framebuffer_init+0x8e7/0xb40 [amdgpu] v2: check adev->enable_virtual_display instead as vkms can be enabled in bare metal as well. CVE-2022-48850 linux - multiple products In the Linux kernel, the following vulnerability has been resolved: net-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| enabled in bare metal as well. CVE-2022-48850 linux - multiple products In the Linux kernel, the following vulnerability has been resolved: 2024-07-16 5.5 Modern to speed_show Net-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| products net-sysfs: add check for netdevice being present to speed_show When bringing down the netdevice or system shutdown, a panic | edium |
| When bringing down the netdevice or system shutdown, a panic | |
| | |
| triggered while accessing the sysfs path because the device is | |
| already removed. | |
| [755.549084] mlx5_core 0000:12:00.1: Shutdown was called [756.404455] mlx5_core 0000:12:00.0: Shutdown was called | |
| [757.937260] BUG: unable to handle kernel NULL pointer dereference at (null) [758.031397] IP: [<fffffff8ee11acb>] dma_pool_alloc+0x1ab/0x280</fffffff8ee11acb> | |
| crash> bt | |
| PID: 12649 TASK: ffff8924108f2100 CPU: 1 COMMAND: "amsd" | |
| #9 [ffff89240e1a38b0] page_fault at fffffff8f38c778 [exception RIP: dma_pool_alloc+0x1ab] RIP: ffffffff8ee11acb RSP: ffff89240e1a3968 RFLAGS: 00010046 | |
| RAX: 0000000000000246 RBX: ffff89243d874100 RCX: 0000000000001000 RDX: 00000000000000 RSI: 00000000000246 RDI: | |
| ffff89243d874090 RBP: ffff89240e1a39c0 R8: 0000000001f080 R9: ffff8905ffc03c00 | |
| R10: fffffffc04680d4 R11: ffffffff8edde9fd R12: 00000000000000080d0 | |
| R13: ffff89243d874090 R14: ffff89243d874080 R15: 000000000000000 ORIG_RAX: ffffffffffff CS: 0010 SS: 0018 | |
| #10 [ffff89240e1a39c8] mlx5_alloc_cmd_msg at ffffffffc04680f3 [mlx5_core] | |
| #11 [ffff89240e1a3a18] cmd_exec at fffffffc046ad62 [mlx5_core] | |
| #12 [ffff89240e1a3ab8] mlx5_cmd_exec at fffffffc046b4fb [mlx5_core] #13 [ffff89240e1a3ae8] mlx5_core_assess_reg_at | |
| #13 [ffff89240e1a3ae8] mlx5_core_access_reg at fffffffc0475434 [mlx5_core] #14 [ffff89240e1a3b40] mlx5e_get_fec_caps at fffffffc04a7348 | |
| [mlx5_core] #15 [ffff89240e1a3bb0] get_fec_supported_advertised at | |
| ffffffffc04992bf [mlx5_core] #16 [ffff89240e1a3c08] mlx5e_get_link_ksettings at | |
| ffffffffc049ab36 [mlx5_core] #17 [ffff89240e1a3ce8]ethtool_get_link_ksettings at fffffffff8f25db46 | |

| | | #18 [ffff89240e1a3d48] speed_show at ffffffff8f277208 #19 [ffff89240e1a3dd8] dev_attr_show at fffffff8f0b70e3 #20 [ffff89240e1a3df8] sysfs_kf_seq_show at fffffff8eedbedf #21 [ffff89240e1a3e18] kernfs_seq_show at fffffff8eeda596 #22 [ffff89240e1a3e28] seq_read at fffffff8ee76d10 #23 [ffff89240e1a3e98] kernfs_fop_read at fffffff8eedaef5 #24 [ffff89240e1a3ed8] vfs_read at fffffff8ee4e3ff #25 [ffff89240e1a3f08] sys_read at fffffff8ee4f27f #26 [ffff89240e1a3f50] system_call_fastpath at ffffffff8f395f92 crash> net_device.state ffff89443b0c0000 state = 0x5 (LINK_STATE_START LINK_STATE_NOCARRIER) | | | |
|-----------------------|------------------|--|------------|-----|---------|
| | | To prevent this scenario, we also make sure that the netdevice is present. | | | |
| CVE-2022-48853 | linux - multiple | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | swiotlb: fix info leak with DMA_FROM_DEVICE | | | |
| | | The problem I'm addressing was discovered by the LTP test covering cve-2018-1000204. | | | |
| | | A short description of what happens follows: 1) The test case issues a command code 00 (TEST UNIT READY) via the SG_IO | | | |
| | | interface with: dxfer_len == 524288, dxdfer_dir == SG_DXFER_FROM_DEV and a corresponding dxferp. The peculiar thing about this is that | | | |
| | | TUR is not reading from the device. | | | |
| | | 2) In sg_start_req() the invocation of blk_rq_map_user() effectively bounces the user-space buffer. As if the device was to transfer | | | |
| | | into it. Since commit a45b599ad808 ("scsi: sg: allocate with | | | |
| | | GFP_ZERO in sg_build_indirect()") we make sure this first bounce buffer is allocated with GFP_ZERO. | | | |
| | | 3) For the rest of the story we keep ignoring that we have a TUR, so the | | | |
| | | device won't touch the buffer we prepare as if the we had a DMA_FROM_DEVICE type of situation. My setup uses a virtio-scsi device | | | |
| | | and the buffer allocated by SG is mapped by the function virtqueue_add_split() which uses DMA_FROM_DEVICE for the "in" sgs (here | | | |
| | | scatter-gather and not scsi generics). This mapping involves bouncing via the swiotlb (we need swiotlb to do virtio in protected guest | | | |
| | | like s390 Secure Execution, or AMD SEV). 4) When the SCSI TUR is done, we first copy back the content of | | | |
| | | the second (that is swiotlb) bounce buffer (which most likely contains some previous IO data), to the first bounce buffer, which contains all zeros. Then we copy back the content of the first bounce buffer | | | |
| | | to the user-space buffer. 5) The test case detects that the buffer, which it zero-initialized, ain't all zeros and fails. | | | |
| | | One can argue that this is an swiotlb problem, because without swiotlb we leak all zeros, and the swiotlb should be transparent in a sense | | | |
| | | that it does not affect the outcome (if all other participants are well behaved). | | | |
| | | Copying the content of the original buffer into the swiotlb buffer is the only way I can think of to make swiotlb transparent in such scenarios. So let's do just that if in doubt, but allow the driver | | | |
| | | to tell us that the whole mapped buffer is going to be overwritten, in which case we can preserve the old behavior and avoid the performance | | | |
| CVE-2022-48856 | linux - multiple | impact of the extra bounce. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| <u>5.1 2522 40000</u> | products | gianfar: ethtool: Fix refcount leak in gfar_get_ts_info | | 3.3 | caraiii |
| | - | · | | • | • |

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|----------------|------------------------------|--|------------|-----|--------|
| CVE-2022-48857 | linux - multiple products | The of_find_compatible_node() function returns a node pointer with refcount incremented, We should use of_node_put() on it when done Add the missing of_node_put() to release the refcount. In the Linux kernel, the following vulnerability has been resolved: NFC: port100: fix use-after-free in port100_send_complete | 2024-07-16 | 5.5 | Medium |
| | | Syzbot reported UAF in port100_send_complete(). The root case is in missing usb_kill_urb() calls on error handling path of ->probe function. | | | |
| | | port100_send_complete() accesses devm allocated memory which will be freed on probe failure. We should kill this urbs before returning an error from probe function to prevent reported use-after-free | | | |
| | | Fail log: | | | |
| | | BUG: KASAN: use-after-free in port100_send_complete+0x16e/0x1a0 drivers/nfc/port100.c:935 Read of size 1 at addr ffff88801bb59540 by task ksoftirqd/2/26 | | | |
| | | Call Trace: <task> dump_stack lib/dump_stack.c:88 [inline] dump_stack_lvl+0xcd/0x134 lib/dump_stack.c:106 print_address_description.constprop.0.cold+0x8d/0x303 mm/kasan/report.c:255 kasan_report mm/kasan/report.c:442 [inline] kasan_report.cold+0x83/0xdf mm/kasan/report.c:459 port100_send_complete+0x16e/0x1a0 drivers/nfc/port100.c:935 usb_hcd_giveback_urb+0x2b0/0x5c0 drivers/usb/core/hcd.c:1670</task> | | | |
| | | | | | |
| | | Allocated by task 1255: kasan_save_stack+0x1e/0x40 mm/kasan/common.c:38 kasan_set_track mm/kasan/common.c:45 [inline] set_alloc_info mm/kasan/common.c:436 [inline]kasan_kmalloc mm/kasan/common.c:515 [inline]kasan_kmalloc mm/kasan/common.c:474 [inline]kasan_kmalloc+0xa6/0xd0 mm/kasan/common.c:524 alloc_dr drivers/base/devres.c:116 [inline] devm_kmalloc+0x96/0x1d0 drivers/base/devres.c:823 devm_kzalloc include/linux/device.h:209 [inline] port100_probe+0x8a/0x1320 drivers/nfc/port100.c:1502 | | | |
| | | Freed by task 1255: kasan_save_stack+0x1e/0x40 mm/kasan/common.c:38 kasan_set_track+0x21/0x30 mm/kasan/common.c:45 kasan_set_free_info+0x20/0x30 mm/kasan/generic.c:370kasan_slab_free mm/kasan/common.c:366 [inline]kasan_slab_free+0xff/0x140 mm/kasan/common.c:328 kasan_slab_free include/linux/kasan.h:236 [inline]cache_free mm/slab.c:3437 [inline] kfree+0xf8/0x2b0 mm/slab.c:3794 release_nodes+0x112/0x1a0 drivers/base/devres.c:501 devres_release_all+0x114/0x190 drivers/base/devres.c:530 | | | |
| CVE-2022-48859 | linux - multiple products | really_probe+0x626/0xcc0 drivers/base/dd.c:670 In the Linux kernel, the following vulnerability has been resolved: net: marvell: prestera: Add missing of_node_put() in prestera_switch_set_base_mac_addr | 2024-07-16 | 5.5 | Medium |
| | | This node pointer is returned by of_find_compatible_node() with refcount incremented. Calling of_node_put() to aovid the refcount leak. | | | |
| CVE-2022-48860 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | | ethernet: Fix error handling in xemaclite_of_probe | | | |
| | | This node pointer is returned by of_parse_phandle() with refcount incremented in this function. Calling of_node_put() to avoid the refcount leak. As the remove function do. | | | |

| CVE-2022-48861 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
|-----------------------|------------------------------|--|-------------------------|-----|------------|
| | products | vdpa: fix use-after-free on vp_vdpa_remove | | | |
| | | When vp_vdpa driver is unbind, vp_vdpa is freed in | | | |
| | | vdpa_unregister_device and then vp_vdpa->mdev.pci_dev is dereferenced in | | | |
| | | vp_modern_remove, | | | |
| | | triggering use-after-free. | | | |
| | | Call Trace of unbinding driver free vp_vdpa: | | | |
| | | do_syscall_64 vfs_write | | | |
| | | kernfs_fop_write_iter | | | |
| | | device_release_driver_internal pci_device_remove | | | |
| | | vp_vdpa_remove | | | |
| | | vdpa_unregister_device kobject_release | | | |
| | | device_release | | | |
| | | kfree | | | |
| | | Call Trace of dereference vp_vdpa->mdev.pci_dev: | | | |
| | | vp_modern_remove pci_release_selected_regions | | | |
| | | pci_release_region | | | |
| | | pci_resource_len pci_resource_end | | | |
| | | (dev)->resource[(bar)].end | | | |
| CVE-2022-48862 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | products | vhost: fix hung thread due to erroneous iotlb entries | | | |
| | | In vhost_iotlb_add_range_ctx(), range size can overflow to 0 when | | | |
| | | start is 0 and last is ULONG_MAX. One instance where it can happen | | | |
| | | is when userspace sends an IOTLB message with iova=size=uaddr=0 | | | |
| | | (vhost_process_iotlb_msg). So, an entry with size = 0, start = 0, | | | |
| | | last = ULONG_MAX ends up in the iotlb. Next time a packet is sent, iotlb_access_ok() loops indefinitely due to that erroneous entry. | | | |
| | | Call Trace: | | | |
| | | <task> iotlb_access_ok+0x21b/0x3e0 drivers/vhost/vhost.c:1340</task> | | | |
| | | vq_meta_prefetch+0xbc/0x280 | | | |
| | | drivers/vhost/vhost.c:1366 vhost_transport_do_send_pkt+0xe0/0xfd0 | | | |
| | | drivers/vhost/vsock.c:104 | | | |
| | | vhost_worker+0x23d/0x3d0 drivers/vhost/vhost.c:372 kthread+0x2e9/0x3a0 kernel/kthread.c:377 | | | |
| | | ret_from_fork+0x1f/0x30 arch/x86/entry/entry_64.S:295 | | | |
| | | | | | |
| | | Reported by syzbot at: | | | |
| | | https://syzkaller.appspot.com/bug?extid=0abd373e2e50d 704db87 | | | |
| | | To fix this, do two things: | | | |
| | | Return -EINVAL in vhost_chr_write_iter() when userspace asks | | | |
| | | to map a range with size 0. | | | |
| | | 2. Fix vhost_iotlb_add_range_ctx() to handle the range [0, ULONG_MAX] | | | |
| CVE-2022-48863 | linux - multiple | by splitting it into two entries. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| CVL-2U22-40003 | products | | ∠∪∠ + -∪/-10 | J.J | ivicululli |
| | | mISDN: Fix memory leak in dsp_pipeline_build() | | | |
| | | dsp_pipeline_build() allocates dup pointer by kstrdup(cfg), | | | |
| | | but then it updates dup variable by strsep(&dup, " "). As a result when it calls kfree(dup), the dup variable contains | | | |
| | | NULL. | | | |
| | | Found by Linux Driver Verification project (linuxtesting.org) with | | | |
| CVE 2022 400C4 | liming magalated a | SVACE. | 2024 07 40 | | N/odi |
| <u>CVE-2022-48864</u> | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | 1 | | <u> </u> | 1 | <u>I</u> |

| 1 | | vdpa/mlx5: add validation for VIRTIO_NET_CTRL_MQ_VQ_PAIRS_SET command | | | |
|----------------|------------------------------|---|------------|-----|--------|
| | | When control vq receives a VIRTIO_NET_CTRL_MQ_VQ_PAIRS_SET | | | |
| | | command request from the driver, presently there is no validation against | | | |
| | | the number of queue pairs to configure, or even if multiqueue had | | | |
| | | been negotiated or not is unverified. This may lead to kernel panic due | | | |
| | | to uninitialized resource for the queues were there any bogus | | | |
| | | request sent down by untrusted driver. Tie up the loose ends there. | | | |
| CVE-2022-48865 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 5.5 | Medium |
| | p. 5 | tipc: fix kernel panic when enabling bearer | | | |
| | | When enabling a bearer on a node, a kernel panic is observed: | | | |
| | | [4.498085] RIP: 0010:tipc_mon_prep+0x4e/0x130 [tipc] | | | |
| | | [4.520030] Call Trace: [4.520689] <irq> [4.521236] tipc_link_build_proto_msg+0x375/0x750 [tipc] [4.522654] tipc_link_build_state_msg+0x48/0xc0 [tipc] [4.524034]tipc_node_link_up+0xd7/0x290 [tipc] [4.525292] tipc_rcv+0x5da/0x730 [tipc] [4.526346] ?netif_receive_skb_core+0xb7/0xfc0 [4.527601] tipc_l2_rcv_msg+0x5e/0x90 [tipc] [4.528737]netif_receive_skb_list_core+0x20b/0x260 [4.530068] netif_receive_skb_list_internal+0x1bf/0x2e0 [4.531450] ? dev_gro_receive+0x4c2/0x680</irq> | | | |
| | | [4.532512] napi_complete_done+0x6f/0x180 | | | |
| | | [4.533570] virtnet_poll+0x29c/0x42e [virtio_net] | | | |
| | | The node in question is receiving activate messages in another thread after changing bearer status to allow message sending/receiving in current thread: | | | |
| | | thread 1 thread 2 | | | |
| | | <pre>tipc_enable_bearer() test_and_set_bit_lock() tipc_bearer_xmit_skb() tipc_l2_rcv_msg() tipc_rcv() tipc_node_link_up() tipc_link_build_state_msg() tipc_link_build_proto_msg() tipc_mon_prep() {</pre> | | | |
| | | // Not being executed yet tipc_mon_create() | | | |
| | | { | | | |
| | | } | | | |
| | | Monitoring pointer in thread 2 is dereferenced before monitoring data | | | |
| | | is allocated in thread 1. This causes kernel panic. | | | |
| | | This commit fixes it by allocating the monitoring data before enabling | | | |
| CVE-2024-41009 | linux - multiple | the bearer to receive messages. In the Linux kernel, the following vulnerability has been resolved: | 2024-07-17 | 5.5 | Medium |
| | products | bpf: Fix overrunning reservations in ringbuf | | | |

The BPF ring buffer internally is implemented as a power-of-2 sized circular buffer, with two logical and ever-increasing counters: consumer_pos is the consumer counter to show which logical position the consumer consumed the data, and producer_pos which is the producer counter denoting the amount of data reserved by all producers. Each time a record is reserved, the producer that "owns" the record will successfully advance producer counter. In user space each time a record is read, the consumer of the data advanced the consumer counter once it finished processing. Both counters are stored in separate pages so that from user space, the producer counter is read-only and the consumer counter is read-write. One aspect that simplifies and thus speeds up the implementation of both producers and consumers is how the data area is mapped twice contiguously back-to-back in the virtual memory, allowing to not take any special measures for samples that have to wrap around at the end of the circular buffer data area, because the next page after the last data page would be first again, and thus the sample will still appear completely contiguous in virtual memory. Each record has a struct bpf_ringbuf_hdr { u32 len; u32 pg_off; } header for book-keeping the length and offset, and is inaccessible to the BPF program. Helpers like bpf_ringbuf_reserve() return `(void *)hdr + BPF_RINGBUF_HDR_SZ` for the BPF program to use. Bing-Jhong and Muhammad reported that it is however possible to make a second allocated memory chunk overlapping with the first chunk and as a result, the BPF program is now able to edit first chunk's header. For example, consider the creation of a BPF_MAP_TYPE_RINGBUF map with size of 0x4000. Next, the consumer_pos is modified to 0x3000 /before/ a call to bpf ringbuf reserve() is made. This will allocate a chunk A, which is in [0x0,0x3008], and the BPF program is able to edit [0x8,0x3008]. Now. lets allocate a chunk B with size 0x3000. This will succeed because consumer_pos was edited ahead of time to pass the 'new_prod_pos - cons_pos > rb->mask` check. Chunk B will be in range [0x3008,0x6010], and the BPF program is able to edit [0x3010,0x6010]. Due to the ring buffer memory layout earlier, the ranges [0x0,0x4000] and [0x4000,0x8000] point to the same data pages. This means that chunk B at [0x4000,0x4008] is chunk A's bpf_ringbuf_submit() / bpf_ringbuf_discard() use the header's pg_off to then locate the bpf_ringbuf itself via bpf_ringbuf_restore_from_rec(). Once chunk

B modified chunk A's header, then bpf ringbuf commit() refers to

Fix it by calculating the oldest pending pos and check whether the

the wrong

page and could cause a crash.

| products bpf: Fix too early release of tox_entry Pedro Pinto and later independently also Hyunwoo Kim and Wong! Lee reported an issue that the tox_entry can be released too early leading to a use after free (UAF) when an active old style ingress or clast t gdisc with a shared to block is later replaced by another ingress or clast instance. Essentially, the sequence to trigger the UAF (one exemple) can be as follows: 1. A network nomespace is created 2. An ingress gdisc is created. This allocates a tox_entry, and &tox_entry>miniq is stored in the qdisc's miniqp >p_miniq. At the same time, a sto block with index 1 is created. 3. chain is statched to the tollobus chaind must be connected to the block linked to the ingress, gdisc to later reach the function tof, chaind, head, change, cb. deli) which triggers the UAF. 4. Create and graft a clast adjoint. This causes the ingress gdisc created in step 1 to be removed, thus freeing the previously linked tox_entry. Intellink_row_msg() >> to, contry, rheetlink_row_msg() >> pdisc_create() >> pdisc_create() >> Finally, the network namespace is closed. This registers the delarup_net worker, and uning the process of releasing the remaining closet adjox, it accesses the tox_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() >> delaru_device_ost_batch() >> delaru_device_ost_batch() >> delaru_device_ost_batch() >> delaru_device_ost_batch() >> delaru_device_ost_batch() >> delar_device_ost_batch() >> dela | | | from the oldest outstanding record to the newest would span beyond the ring buffer size. If that is the case, then reject the request. We've tested with the ring buffer benchmark in BPF selftests (./benchs/run_bench_ringbufs.sh) before/after the fix and while it seems a bit slower on some benchmarks, it is still not significantly enough to matter. | | | |
|--|----------------|----------|---|------------|-----|--------|
| Pedro Pinto and later independently also Hyunwoo Kim and Wongl Lee reported an issue that the toc_entry can be released too early leading to a use after free (UAF) when an active old-style ingress or cloact glidsc with a shared to block is later replaced by another ingress or cloact instance. Essentially, the sequence to trigger the UAF (one example) can be as follows: 1. A network namespace is created 2. An ingress galoc is created. This allocates a tox_entry, and &tox_entry>minql is stored in the galoc's minqp>p_minq. At the same time, a tell block with index 1 is created. 3. A chang is attached to the tell block. Chain of must be connected to the block linked to the ingress galoc to later reach the function tot_chain(). head_change_to_dell) which triggers the UAF. 4. Create and graff a cloate glote. This causes the ingress galoc created in step 1 to be removed, thus freeing the previously linked tox_entry: thetlink_row_msg[>> tu_modif_galoc() >> tugisc_create() >> tugisc_create(| CVE-2024-41010 | <u> </u> | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-17 | 5.5 | Medium |
| Lee reported an issue that the txc_entry can be released too early leading to a use after free (UAF) when an active old-style ingress or clact cipiss with a shared to block is later replaced by another ingress or chact instance. Essentially, the sequence to trigger the UAF (one example) can be as follows: 1. A network namespace is created 2. An ingress gdisc is created. This allocates a txc_entry, and &txc_entry-ariniq is strong in the edisc's minign>p_miniq. At the same time, a tcf block with index 1 is created. 3. Chain 0 is attached to the tcf block. Chain 0 must be connected to the block linked to the ingress gdisc to later reach the function tcf_chain 0. head_change_tcd_del() which triggers the UAF. 4. Create and graft a cload tolic. This causes the ingress gdisc created in step 1 to be removed, thus freeing the previously linked txc_entry: truetlink_rcv_msg() => tc_modify_gdisc() => clastint[(a)] => qdisc_clastor() => clastint[(a)] => qdisc_clastor() => clastint[(a)] => pdisc_clastor() => pdisc_destor() => pdisc_destor() => pdisc_destor() => pdisc_destor() => pdisc_destor() => pdisc_destor() => point_int() => unregister_netdevice_many() => unregister_netdevice_many() => pdisc_destor() => clastint[(a)] => pdisc_destor() => clastint() => cla | | | bpf: Fix too early release of tcx_entry | | | |
| an issue that the toc_entry can be released too early leading to a use after free (UAF) when an active old-style ingress or cloact gdisc with a shared to block is later replaced by another ingress or cloact instance. Essentially, the sequence to trigger the UAF (one example) can be as follows: 1. A network namespace is created 2. An ingress gdisc is created. This allocates a tox_entry, and 8 tox_entry-miniq as stored in the gdisc's miniqs >p_miniq. At the same time, a tcf block with index 1 is created. 3. Chando is attached to the tcf block, chaind most be connected to the block linked to the ingress gdisc to later each the function tof, chaino, head, change to, adely which ringges the UAF. 4. Create and graft a cloact todisc. This causes the ingress optics created in step 1 to be removed, thus freeing the previously linked tox_entry: **thetime, rev_msg()** ⇒> (z_modify_gdisc()** ⇒> disc_created()** >> close_created()** >> close_created(| | | | | | |
| after free (UAF) when an active old style ingress or chact gdisc with a shared to block is later replaced by another ingress or clasact instance. Essentially, the sequence to trigger the UAF (one example) can be as follows: 1. A network namespace is created 2. An ingress addisc is created. This allocates a tox, entry, and 8tcx, entry-prinking is stored in the pdisc's minigp->p_miniq. At the same time, a tof block with index 1 is created. 3. chalno is attached to the tof block, chaind must be connected to the block linked to the lingress addisc to later reach the function tof chaino. head, change, cb. delly which triggers the UAF. 4. Create and graft a clasted right. This causes the ingress addisc created in step 1 to be removed, thus freeing the previously linked tot, entry: Intellink_rev_msg() >> tc_modify_qdisc() >> clast_mint() [a] >> daisc_graft() >> clast_grift() >> daisc_graft() >> firee_rev() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup, net worker, and during the process of releasing the remaining clast_disk_t. accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() >> ops_cet_list() >> default_device_exit_batch() >> unregister_netdevice_many() >> unregister_netdevice_many() >> unregister_netdevice_many() >> tof_chain_head_change() >> tof_chain_head_change() >> inini_qdisc_poir_swap()// UAF There are also other variants, the gist is to add an ingress (or close() qdisc_with a specific shared block, then to replace that qdisc, waiting | | | an issue that the tcx_entry can be released too early leading to a | | | |
| shared to block is later replaced by another ingress or clsact instance. Essentially, the sequence to trigger the UAF (one example) can be as follows: 1. A network namespace is created. 2. An ingress golds is created. This allocates a tox_entry, and &tox_entry_sminiq is stored in the qdsx's miniqp>>p_miniq. At the same time, a tof block with index 1 is created. 3. chain'd is attached to the tof block. chain'd must be connected to the block linked to the ingress gdisc to later reach the function tof_chaino_head_change_cb_del[) which triggers the UAF. 4. Create and graft a Casted ordisc. This causes the ingress gdisc created in step 1 to be removed, thus freeing the previously linked tot_entry: **rtnetlink_rcv_msg()** **> chact_intil() [a]** *> palisc_gestroy() [b]** *>> calisc_gestroy() [b]** *> qalisc_gestroy() [b]** *> larges_gestry() [to]** *> unregister_gestry() [to]** *> larges_gestry() [to, larges_gestry() [t | | | after free (UAF) when an active old-style ingress or clsact qdisc | | | |
| as follows: 1. A network namespace is created 2. An ingress adisc is created. This allocates a tox_entry, and 8.tox_entry->miniq is stored in the qdisc's miniqp>p_miniq.At the same time, a tof block with index 1 is created. 3. chain 0 is attached to the tof block, chain 0 must be connected to the block linked to the ingress qdisc to later reach the function tof chain 0, head, change_cb_del() which triggers the UAF. 4. Create and graft a cleart qdisc. This causes the ingress qdisc created in step 1 to be removed, thus freeing the previously linked tox_entry: rtnetlink_rcv_msg() >> tc_modify_qdisc() >> qdisc_create() >> clear_clear | | | shared tc block is later replaced by another ingress or clsact | | | |
| 2. An ingress gdisc is created. This allocates at cx_entry, and & & tx_entry-miniq is stored in the qdisc's miniqp->p_miniq. At the same time, a tcf block with index 1 is created. 3. chain0 is attached to the tcf block, chain0 must be connected to the block linked to the ingress qdisc to later reach the function tcf_chain0_head_change_cb_del() which triggers the UAF. 4. Create and graft a clast qdisc. This causes the ingress qdisc created in step 1 to be removed, thus freeling the previously linked tox_entry: trinetlink_rcv_msg() >> tc_modify_qdisc() >> tc_modify_qdisc() >> qdisc_create() >> clast_init() [a] >> qdisc_destroy() >> ingress_destroy() >> ingress_destroy() >> ingress_destroy() >> ingress_destroy() >> ingress_destroy() [b] >> tc_modify_fel() >> kfree_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clast qdisc, it accesses the tx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() >> op_ge_exit_list() >> op_ge_exit_list() >> op_ge_exit_list() >> op_ge_exit_list() >> op_ge_exit_list() >> clast_device_exit_batch() >> tcf_chain0_head_change_tem() >> tcf_chain0_head_change() >> tcf_chain_head_change() >> min_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clast()) qdisc with a specific shared block, then to replace that qdisc, waiting | | | , | | | |
| same time, a tcf block with index 1 is created. 3 chain0 is attached to the tcf block, chain0 must be connected to the block linked to the ingress gdisc to later reach the function tcf_chain0_head_change_cb_del() which triggers the UAF. 4. Create and graft a clsact qdisc. This causes the lingress gdisc created in step 1 to be removed, thus freeing the previously linked tcx_entry: rtnetlink_rcv_msg() >> tc_modify_qdisc() => clsact_init() [a] >> qdisc_graft() >> qdisc_graft() >> lessed_setroy() >> lessed_setroy() [b] >> tcx_entry_free() >> kfree_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() >> ops_exit_isit() >> ops_exit_isit() >> unregister_netdevice_many() >> unregister_netdevice_many() >> unregister_netdevice_many() >> dsact_destroy() [c] >> tcf_olock_put_exit_olock_p | | | 2. An ingress qdisc is created. This allocates a tcx_entry, and &tcx_entry->miniq is stored in the qdisc's miniqp->p_miniq. At | | | |
| the block linked to the ingress adisc to later reach the function tof_chainO_head_change_cb_del() which triggers the UAF. 4. Create and graft a clsact adisc. This causes the ingress adisc created in step 1 to be removed, thus freeing the previously linked tox_entry: rtnetlink_rcv_msg() >> tc_modify_adisc() >> qdisc_create() >> clsact_init() [a] >> qdisc_gestroy() >> lags_graft() >> tcx_entry_free() >> tcx_entry_free() >> tcx_entry_free() >> free_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact adisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() >> ops_exit_list() => default_device_exit_batch() >> unregister_netdevice_many_notlify() >> des_vshutdown() >> des_vshutdown() >> des_vshutdown() >> clsact_destroy() [c] >> tcf_chain_head_change_cb_del() >> tcf_chain_head_change_() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc_with a specific shared block, then to replace that qdisc, waiting | | | same time, a tcf block with index 1 is created. 3. chain0 is attached to the tcf block. chain0 must be connected | | | |
| rtnetlink_rcv_msg() => tc_modify_qdisc() => qdisc_create() => clasat_init() [a] => qdisc_graft() => qdisc_destroy() => ingress_destroy() [b] => tcx_entry_free() => kfree_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_ext_list() => ops_ext_list() => unregister_netdevice_many() => unregister_netdevice_many_notify() => desu_shutdown() => disc_put() => clsact_destroy() [c] => tf_chain_head_change_cb_del() => tf_chain_head_change() => min_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc_with a specific shared block, then to replace that qdisc, waiting | | | the block linked to the ingress qdisc to later reach the function tcf_chainO_head_change_cb_del() which triggers the UAF. 4. Create and graft a clsact qdisc. This causes the ingress qdisc created in step 1 to be removed, thus freeing the previously | | | |
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| => clsact_init() [a] => qdisc_garfit() => qdisc_destroy() =>qdisc_destroy() => ingress_destroy() [b] => txentry_free() => kfree_rcu() // txentry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the txentry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => tcf_block_put_ext() => tcf_chaino_head_change_to_del() => tcf_chaino_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| => qdisc_destroy() => qdisc_destroy() => largess_destroy() [b] => tox_entry_free() => kfree_rcu() // tox_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tox_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many() => largester_netdevice_many_notify() => dev_shutdown() => dev_shutdown() => dev_shutdown() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chainO_head_change_cb_del() => tcf_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => qdisc_create() | | | |
| =>qdisc_destroy() => ingress_destroy() [b] => tox_cntry_free() => kfree_rcu() // tox_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tox_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => ofs_exit_list() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tof_block_put_exit() => tof_chain_head_change_them() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => qdisc_graft() | | | |
| => tcx_entry_free() => kfree_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chain_head_change_cb_del() => tcf_chain_head_change() => tcf_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| => kfree_rcu() // tcx_entry freed 5. Finally, the network namespace is closed. This registers the cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chain_head_change_cb_del() => tcf_chain_head_change_() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was already freed in step 4, causing the UAF to occur: cleanup_net() => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chaino_head_change_cb_del() => tcf_chain_head_change() => clsact_destroy() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| => ops_exit_list() => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chainO_head_change_cb_del() => tcf_chain_head_change() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | cleanup_net worker, and during the process of releasing the remaining clsact qdisc, it accesses the tcx_entry that was | | | |
| => default_device_exit_batch() => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chainO_head_change_cb_del() => tcf_chain head_change() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | cleanup_net() | | | |
| => unregister_netdevice_many() => unregister_netdevice_many_notify() => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chain0_head_change_cb_del() => tcf_chain_head_change_item() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| => dev_shutdown() => qdisc_put() => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chain_head_change_cb_del() => tcf_chain_head_change_item() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => unregister_netdevice_many() | | | |
| => clsact_destroy() [c] => tcf_block_put_ext() => tcf_chain0_head_change_cb_del() => tcf_chain_head_change_item() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => dev_shutdown() | | | |
| => tcf_block_put_ext() => tcf_chain0_head_change_cb_del() => tcf_chain_head_change_item() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| => tcf_chain_head_change_item() => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => tcf_block_put_ext() | | | |
| => clsact_chain_head_change() => mini_qdisc_pair_swap() // UAF There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| There are also other variants, the gist is to add an ingress (or clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | => clsact_chain_head_change() | | | |
| clsact) qdisc with a specific shared block, then to replace that qdisc, waiting | | | | | | |
| waiting | | | clsact) | | | |
| for the text entry lefter result to be executed and subsequently | | | waiting | | | |
| for the tcx_entry kfree_rcu() to be executed and subsequently accessing | | | accessing | | | |
| the current active qdisc's miniq one way or another. | | | | | | |
| The correct fix is to turn the miniq_active boolean into a counter. What | | | <u>-</u> | | | |

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|----------------|--------------------------------------|---|------------|-----|------------|
| | | can be observed, at step 2 above, the counter transitions from 0->1, at step [a] from 1->2 (in order for the miniq object to remain active | | | |
| | | during the replacement), then in [b] from 2->1 and finally [c] 1->0 with the | | | |
| | | eventual release. The reference counter in general ranges from [0,2] and | | | |
| | | it does not need to be atomic since all access to the counter is protected by the rtnl mutex. With this in place, there is no longer a UAF | | | |
| | | happening and the tcx_entry is freed at the correct time. | | | |
| CVE-2024-39728 | ibm - multiple products | IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 is vulnerable to stored cross-site scripting. This vulnerability allows users to embed arbitrary JavaScript code in the Web UI thus altering the intended functionality potentially leading to | 2024-07-15 | 5.4 | Medium |
| CVE-2024-39735 | ibm - multiple | credentials disclosure within a trusted session. IBM X-Force ID: 295967. IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 is | 2024-07-15 | 5.4 | Medium |
| CVE-2024-39733 | products | vulnerable to cross-site scripting. This vulnerability allows an authenticated user to embed arbitrary JavaScript code in the Web UI thus altering the intended functionality potentially leading to credentials disclosure within a trusted session. IBM X-Force ID: 296002. | 2024-07-13 | 3.4 | iviedidiii |
| CVE-2024-39863 | apache - airflow | Apache Airflow versions before 2.9.3 have a vulnerability that allows an authenticated attacker to inject a malicious link when installing a provider. Users are recommended to upgrade to version 2.9.3, which fixes this issue. | 2024-07-17 | 5.4 | Medium |
| CVE-2024-39737 | ibm - multiple products | IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 could allow a remote attacker to obtain sensitive information when a detailed technical error message is returned in the browser. This information could be used in further attacks against the system. IBM X-Force ID: 296004. | 2024-07-15 | 5.3 | Medium |
| CVE-2024-39740 | ibm - multiple products | IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 displays version information in HTTP requests that could allow an attacker to gather information for future attacks against the system. IBM X-Force ID: 296009. | 2024-07-15 | 5.3 | Medium |
| CVE-2024-39741 | ibm - multiple products | IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 could allow a remote attacker to traverse directories on the system. An attacker could send a specially crafted URL request containing "dot dot" sequences (//) to view arbitrary files on the system. IBM X-Force ID: 296010. | 2024-07-15 | 5.3 | Medium |
| CVE-2024-21176 | oracle - mysql_server | Vulnerability in the MySQL Server product of Oracle MySQL (component: Server: Thread Pooling). Supported versions that are affected are 8.4.0 and prior. Difficult to exploit vulnerability allows low privileged attacker with network access via multiple protocols to compromise MySQL Server. Successful attacks of this vulnerability can result in unauthorized ability to cause a hang or frequently repeatable crash (complete DOS) of MySQL Server. CVSS 3.1 Base Score 5.3 (Availability impacts). CVSS Vector: | 2024-07-16 | 5.3 | Medium |
| CVE-2024-6535 | redhat - service_interconne ct | (CVSS:3.1/AV:N/AC:H/PR:L/UI:N/S:U/C:N/I:N/A:H). A flaw was found in Skupper. When Skupper is initialized with the console-enabled and with console-auth set to Openshift, it configures the openshift oauth-proxy with a static cookie-secret. In certain circumstances, this may allow an attacker to bypass authentication to the Skupper console via a specially-crafted cookie. | 2024-07-17 | 5.3 | Medium |
| CVE-2024-21179 | oracle - multiple products | Vulnerability in the MySQL Server product of Oracle MySQL (component: InnoDB). Supported versions that are affected are 8.0.37 and prior and 8.4.0 and prior. Easily exploitable vulnerability allows high privileged attacker with network access via multiple protocols to compromise MySQL Server. Successful attacks of this vulnerability can result in unauthorized ability to cause a hang or frequently repeatable crash (complete DOS) of MySQL Server. CVSS 3.1 Base Score 4.9 (Availability impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:H/UI:N/S:U/C:N/I:N/A:H). | 2024-07-16 | 4.9 | Medium |
| CVE-2024-21185 | oracle - multiple products | Vulnerability in the MySQL Server product of Oracle MySQL (component: InnoDB). Supported versions that are affected are 8.0.38, 8.4.1 and 9.0.0. Easily exploitable vulnerability allows high privileged attacker with network access via multiple protocols to compromise MySQL Server. Successful attacks of this vulnerability can result in unauthorized ability to cause a hang or frequently repeatable crash (complete DOS) of MySQL Server. CVSS 3.1 Base Score 4.9 (Availability impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:H/UI:N/S:U/C:N/I:N/A:H). | 2024-07-16 | 4.9 | Medium |
| CVE-2022-48842 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: | 2024-07-16 | 4.7 | Medium |

```
ice: Fix race condition during interface enslave
Commit 5dbbbd01cbba83 ("ice: Avoid RTNL lock when re-creating
auxiliary device") changes a process of re-creation of aux device
so ice_plug_aux_dev() is called from ice_service_task() context.
This unfortunately opens a race window that can result in dead-
lock
when interface has left LAG and immediately enters LAG again.
Reproducer:
#!/bin/sh
ip link add lag0 type bond mode 1 milmon 100
ip link set lag0
for n in {1..10}; do
   echo Cycle: $n
   ip link set ens7f0 master lag0
   sleep 1
    ip link set ens7f0 nomaster
done
This results in:
[20976.208697] Workqueue: ice ice_service_task [ice]
[20976.213422] Call Trace:
[20976.215871] __schedule+0x2d1/0x830
[20976.219364] schedule+0x35/0xa0
[20976.222510] schedule_preempt_disabled+0xa/0x10
[20976.227043] mutex lock.isra.7+0x310/0x420
[20976.235071] enum all gids of dev cb+0x1c/0x100 [ib core]
[20976.251215] ib enum roce netdev+0xa4/0xe0 [ib core]
[20976.256192] ib cache setup one+0x33/0xa0 [ib core]
[20976.261079] ib_register_device+0x40d/0x580 [ib_core]
[20976.266139] irdma_ib_register_device+0x129/0x250 [irdma]
[20976.281409] irdma_probe+0x2c1/0x360 [irdma]
[20976.285691] auxiliary_bus_probe+0x45/0x70
[20976.289790] really_probe+0x1f2/0x480
[20976.298509] driver_probe_device+0x49/0xc0
[20976.302609] bus_for_each_drv+0x79/0xc0
[20976.306448] __device_attach+0xdc/0x160
[20976.310286] bus_probe_device+0x9d/0xb0
[20976.314128] device_add+0x43c/0x890
[20976.321287] __auxiliary_device_add+0x43/0x60
[20976.325644] ice_plug_aux_dev+0xb2/0x100 [ice]
[20976.330109] ice_service_task+0xd0c/0xed0 [ice]
[20976.342591] process_one_work+0x1a7/0x360
[20976.350536] worker_thread+0x30/0x390
[20976.358128] kthread+0x10a/0x120
[20976.365547] ret_from_fork+0x1f/0x40
[20976.438030] task:ip
                            state:D stack: 0 pid:213658
ppid:213627 flags:0x00004084
[20976.446469] Call Trace:
[20976.448921] __schedule+0x2d1/0x830
[20976.452414] schedule+0x35/0xa0
[20976.455559] schedule_preempt_disabled+0xa/0x10
[20976.460090] __mutex_lock.isra.7+0x310/0x420
[20976.464364] device_del+0x36/0x3c0
[20976.467772] ice_unplug_aux_dev+0x1a/0x40 [ice]
[20976.472313] ice_lag_event_handler+0x2a2/0x520 [ice]
[20976.477288] notifier_call_chain+0x47/0x70
[20976.481386] __netdev_upper_dev_link+0x18b/0x280
[20976.489845] bond_enslave+0xe05/0x1790 [bonding]
[20976.494475] do_setlink+0x336/0xf50
[20976.502517] __rtnl_newlink+0x529/0x8b0
[20976.543441] rtnl_newlink+0x43/0x60
[20976.546934] rtnetlink_rcv_msg+0x2b1/0x360
[20976.559238] netlink_rcv_skb+0x4c/0x120
[20976.563079] netlink_unicast+0x196/0x230
[20976.567005] netlink_sendmsg+0x204/0x3d0
[20976.570930] sock_sendmsg+0x4c/0x50
[20976.574423] ____sys_sendmsg+0x1eb/0x250
[20976.586807] ___sys_sendmsg+0x7c/0xc0
[20976.606353] \ \_\_sys\_sendmsg+0x57/0xa0
[20976.609930] do_syscall_64+0x5b/0x1a0
[20976.613598] entry_SYSCALL_64_after_hwframe+0x65/0xca
```

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|--|----------------------------|---|------------|-----|------------------|
| | | Command 'ip link set nomaster' causes that ice_plug_aux_dev() is called from ice_service_task() context, aux device is created and associated device->lock is taken. Command 'ip link set master' calls ice's notifier under RTNL lock and that notifier calls ice_unplug_aux_dev(). That function tries to take aux device->lock but this is already taken by ice_plug_aux_dev() in step 1 Later ice_plug_aux_dev() tries to take RTNL lock but this is already taken in step 2 Dead-lock The patch fixes this issue by following changes: Bit ICE_FLAG_PLUG_AUX_DEV is kept to be set during ice_plug_aux_dev() call in ice_service_task() The bit is checked in ice_clear_rdma_cap() and only if it is not set then ice_unplug_aux_dev() is called. If it is set (in other words plugging of aux device was requested and ice_plug_aux_dev() is potentially running) then the function only clears thetruncated | | | |
| CVE-2023-52291 | apache - streampark | In streampark, the project module integrates Maven's compilation capabilities. The input parameter validation is not strict, allowing attackers to insert commands for remote command execution, The prerequisite for a successful attack is that the user needs to log in to the streampark system and have system-level permissions. Generally, only users of that system have the authorization to log in, and users would not manually input a dangerous operation command. Therefore, the risk level of this vulnerability is very low. Background: In the "Project" module, the maven build args "<" operator causes command injection. e.g: "< (curl http://xxx.com)" will be executed as a command injection, Mitigation: all users should upgrade to 2.1.4, The "<" operator will blocked. | 2024-07-17 | 4.7 | Medium |
| CVE-2024-29737 | apache - streampark | In streampark, the project module integrates Maven's compilation capabilities. The input parameter validation is not strict, allowing attackers to insert commands for remote command execution, The prerequisite for a successful attack is that the user needs to log in to the streampark system and have system-level permissions. Generally, only users of that system have the authorization to log in, and users would not manually input a dangerous operation command. Therefore, the risk level of this vulnerability is very low. Mitigation: all users should upgrade to 2.1.4 Background info: Log in to Streampark using the default username (e.g. test1, test2, test3) and the default password (streampark). Navigate to the Project module, then add a new project. Enter the git repository address of the project and input `touch /tmp/success_2.1.2` as the "Build Argument". Note that there is no verification and interception of the special character "`". As a result, you will find | 2024-07-17 | 4.7 | Medium |
| | | that this injection command will be successfully executed after executing the build. In the latest version, the special symbol `is intercepted. | | | |
| CVE-2024-39739 | ibm - multiple products | executing the build. | 2024-07-15 | 4.3 | Medium |
| CVE-2024-39739 CVE-2024-39729 CVE-2024-31979 | • | executing the build. In the latest version, the special symbol `is intercepted. IBM Datacap Navigator 9.1.5, 9.1.6, 9.1.7, 9.1.8, and 9.1.9 is vulnerable to server-side request forgery (SSRF). This may allow an authenticated attacker to send unauthorized requests from the system, potentially leading to network enumeration or facilitating | 2024-07-15 | 4.3 | Medium Medium |

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|----------------|-------------------------------|--|------------|-----|--------|
| | | Previously, StreamPipes allowed users to configure custom endpoints from which to install additional pipeline elements. These endpoints were not properly validated, allowing an attacker to get StreamPipes to send an HTTP GET request to an arbitrary address. This issue affects Apache StreamPipes: through 0.93.0. Users are recommended to upgrade to version 0.95.0, which fixes the issue. | | | |
| CVE-2024-21180 | oracle - multiple products | Vulnerability in the PeopleSoft Enterprise PeopleTools product of Oracle PeopleSoft (component: OpenSearch Dashboards). Supported versions that are affected are 8.59, 8.60 and 8.61. Easily exploitable vulnerability allows low privileged attacker with network access via HTTP to compromise PeopleSoft Enterprise PeopleTools. Successful attacks require human interaction from a person other than the attacker and while the vulnerability is in PeopleSoft Enterprise PeopleTools, attacks may significantly impact additional products (scope change). Successful attacks of this vulnerability can result in unauthorized read access to a subset of PeopleSoft Enterprise PeopleTools accessible data. CVSS 3.1 Base Score 4.1 (Confidentiality impacts). CVSS Vector: (CVSS:3.1/AV:N/AC:L/PR:L/UI:R/S:C/C:L/I:N/A:N). | 2024-07-16 | 4.1 | Medium |
| CVE-2024-30471 | apache - streampipes | Time-of-check Time-of-use (TOCTOU) Race Condition vulnerability in Apache StreamPipes in user self-registration. This allows an attacker to potentially request the creation of multiple accounts with the same email address until the email address is registered, creating many identical users and corrupting StreamPipe's user management. This issue affects Apache StreamPipes: through 0.93.0. Users are recommended to upgrade to version 0.95.0, which fixes the issue. | 2024-07-17 | 3.7 | Low |
| CVE-2024-41007 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: tcp: avoid too many retransmit packets If a TCP socket is using TCP_USER_TIMEOUT, and the other peer retracted its window to zero, tcp_retransmit_timer() can retransmit a packet every two jiffies (2 ms for HZ=1000), for about 4 minutes after TCP_USER_TIMEOUT has 'expired'. The fix is to make sure tcp_rtx_probe0_timed_out() takes icsk->icsk_user_timeout into account. Before blamed commit, the socket would not timeout after icsk->icsk_user_timeout, but would use standard exponential backoff for the retransmits. Also worth noting that before commit e89688e3e978 ("net: tcp: fix unexcepted socket die when snd_wnd is 0"), the issue would last 2 minutes instead of 4. | 2024-07-15 | 3.3 | Low |
| CVE-2022-48852 | linux - multiple products | In the Linux kernel, the following vulnerability has been resolved: drm/vc4: hdmi: Unregister codec device on unbind On bind we will register the HDMI codec device but we don't unregister it on unbind, leading to a device leakage. Unregister our device at unbind. | 2024-07-16 | 3.3 | Low |

Where NCA provides the vulnerability information as published by NIST's وإذ تبقى .NIST's NVD. In addition, it is the entity's or individual's responsibility to ensure the implementation of appropriate recommendations.